

# Dr. Luis A. Moncayo-Martínez

## Optimisation of Manufacturing and Logistics Systems

### ■ Research Topics

- Simulation of manufacturing flexible systems.
- Optimisation of inventory systems.
- Mathematical modelling and computational implementation.
- Parallel implementation of meta-heuristics.

### ■ Research Projects

- Optimizing the fiber optic deployment in rural communities.
- Fast simulation of manufacturing system using R and C++.
- Optimisation of service level of bike—sharing systems (PN2015-1234).

### ■ Publications

- <https://orcid.org/0000-0003-4619-3808>
- <https://scholar.google.com/citations?user=weu876YAAAAJ&hl=en>
- <https://www.researchgate.net/profile/Luis-A-Moncayo-Martinez>

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$$\min TSCC = \zeta \sum_{i=1}^N (\mu_i \sum_{j=1}^{N_i} c_{ij} y_{ij}) + h \sum_{i=1}^N (K_i z_i \sigma_i \sqrt{\zeta_i + \sum_{j=1}^{N_i} t_{ij} y_{ij} - \phi_i}) + h \sum_{i=1}^N (W_i \mu_i \sum_{j=1}^{N_i} t_{ij} y_{ij})$$

$$\min LT = \frac{1}{|D|} \sum_{i \in D} LT_i$$

Subject to:

$$LT_i = \sum_{j=1}^{N_i} t_{ij} y_{ij} + \max_{i:(f,i) \in E} (LT_i) \quad \text{for } i = 1, \dots, N$$

$$\sum_{j=1}^{N_i} y_{ij} = 1 \quad \text{for } i = 1, \dots, N$$

$$\zeta_i + \sum_{j=1}^{N_i} t_{ij} y_{ij} - \phi_i \geq 0 \quad \text{for } i = 1, \dots, N$$

$$\phi_f \leq \zeta_i \quad \text{for } i = 1, \dots, N \text{ if } (f, i)$$

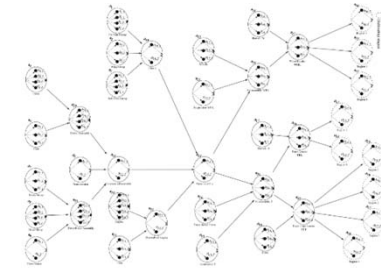
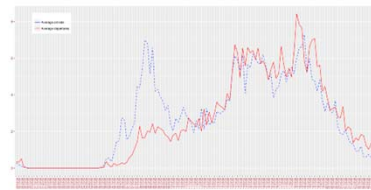
$$\phi_i \leq \Phi_i \quad \text{for all } i \in D$$

$$K_i = \sum_{j=1}^{N_i} c_{ij} y_{ij} + \sum_{(f,i) \in E} K_f$$

$$W_i = \frac{1}{2} \sum_{j=1}^{N_i} c_{ij} y_{ij} + \sum_{(f,i) \in E} K_f$$

$$\phi_i, \zeta_i \geq 0 \quad \text{and integer for } i = 1, \dots, N$$

$$y_{ij} = (0,1) \quad \text{for } i = 1, \dots, N, \quad j \in i, \quad j = 1, \dots, N_i$$



**Algorithm 7** Calculate a solution set  $P$

Require:  $G = (V, E), R > 0, D > 0$

Ensure:  $P = \{s_1, \dots, s_M, \dots\}$

Set  $r = 1$

Initialize  $s_0, v_0, t_0$

If  $r \leq R$  then

  Create the  $r^{\text{th}}$  river with  $D$  drops

  Set  $d = 1$

  while  $d \leq D$  do

*(first part of the algorithm, SC configuration)*

    for  $i = 1$  to  $i = I$  do

      Compute the probability  $p_{ij} \forall j \in i$ , see (7.12)

      Select one option  $j$  based on  $p_{ij}$ , thus (7.4) is solved

      Solve constraints (7.5) and (7.6), thus  $T_i$  and  $C_i$  are set

      Update drop velocity  $v_{d,i}$ , see (7.13)

      Compute soil increments,  $\Delta X_{d,i}$ , see (7.14)

      Update soil of the selected option,  $j \in i$ , see (7.15)

    end for

*(second part of the algorithm, place safety stock)*

    Compute safety stock cost to solve constraints (7.7)–(7.9)

    Calculate the in-transit inventory

    Calculate the  $LT_i$  for every stage  $i$ , see (7.3)

    Compute the solution  $s_d = (LT, IC)$ , see (7.1)–(7.2)

$d++$

  end while

  for all  $s_d$  do

    If  $s_d$  is non-dominated then

$P \leftarrow s_d$

    end if

  end for

  Update soil on selected options  $j$  used to compute  $s_d \in P$ , see (7.16)

$r++$

else

$P_B = P = \{s_1, \dots, s_M, \dots\}$

end if